

Optimal Private Payoff Manipulation against Commitment in Extensive-form games

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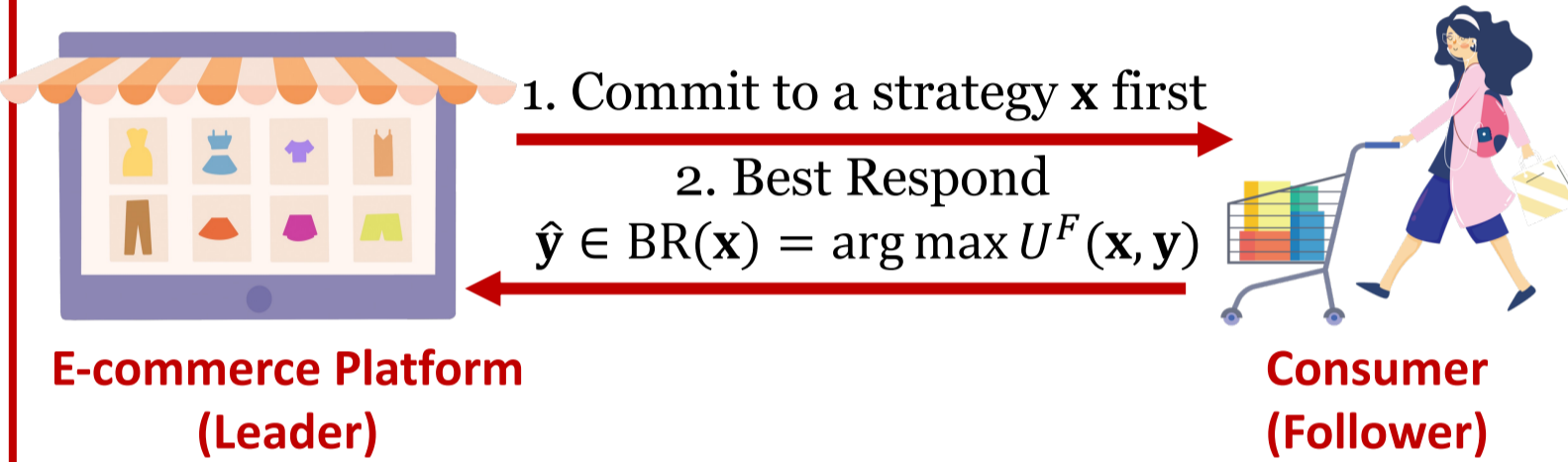


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1. Introduction

Two players play a game: Leader and Follower.



Leader can commit to the optimal strategy maximizing her own utility

$$\hat{x} = \arg \max_{x \in A^L, \hat{y} \in BR(x)} U^L(x, \hat{y})$$

Strong Stackelberg Equilibrium (SSE)

- Leader has the first-mover advantage.
- Follower's payoff information U^F is needed.

Follower can provide **fake** payoff information \tilde{U}^F to

- Get higher utility in the equilibrium of the new game

Practical example: customers of Uber may switch credit cards to the price charged by the platform^[1].

Question here:

How to find the optimal fake \tilde{U}^F that maximizes Follower's actual utility?

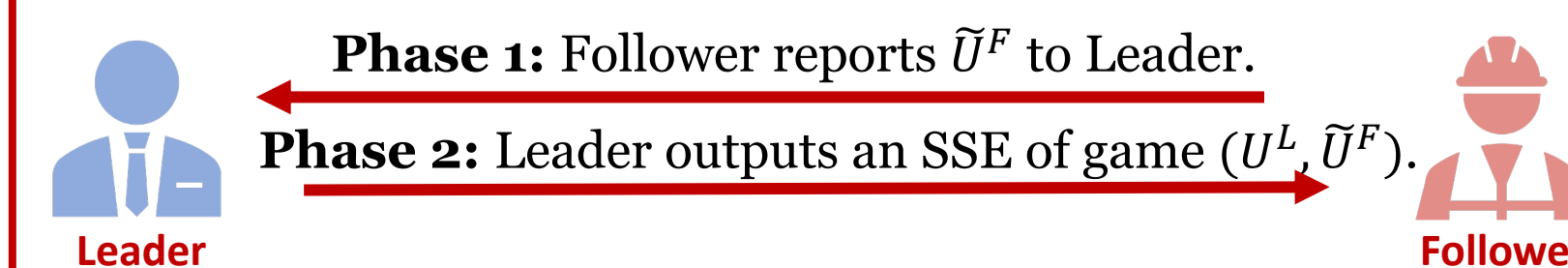
Previous Work:

Focus on **normal-form games**^[2] and **security games**^[3].

Our Work:

Extensive-form games with perfect information.

Model: Two-Phase Game



2. Settings

- SSE with
 - Pure Commitment: A^L the set of all **pure** strategies.
 - Behavioral Commitment: A^L the set of all **behavioral** strategies.
- Probability Distribution p over leaf nodes of the game tree:
 - The **game outcome** the follower aims to finally induce.
- When there exists \tilde{U}^F , such that
 - Inducibility:** One SSE of game (U^L, \tilde{U}^F) leads to p
 - Strong Inducibility:** All SSEs of game (U^L, \tilde{U}^F) lead to p

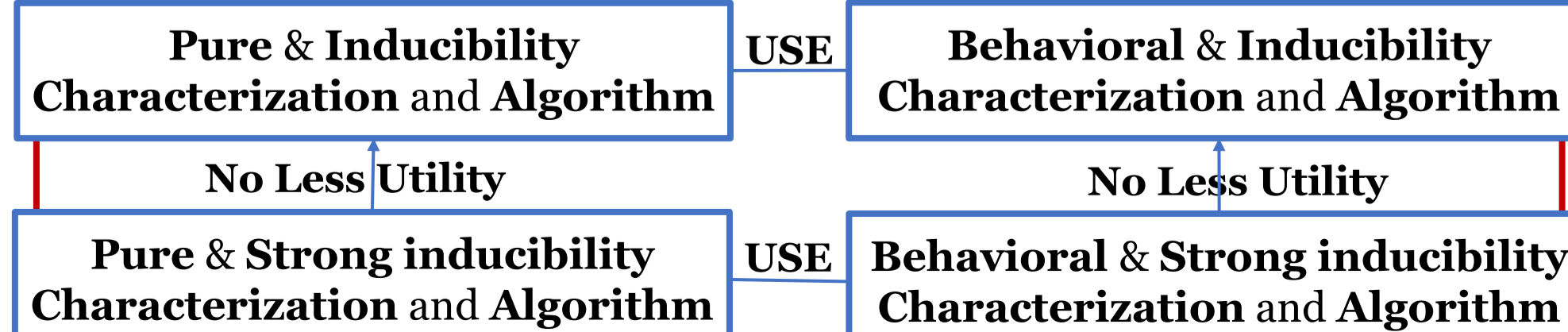
3. Our Contribution

For both pure and behavioral commitment settings:

- Characterizations** of all the (strongly) inducible distributions;
- Polynomial-time Algorithms** for the follower to find an (near-) optimal distribution among all the (strongly) inducible ones, and construct a corresponding follower's payoff function that induces this distribution.

WE compare the **optimal** utilities in **one** game that a follower can get in the four different settings:

- Inducibility v.s. Strong Inducibility:** characterization of the games where the two values are (nearly-)equal.
 - Utility Supremum Equivalence (USE) property
- Pure v.s. Behavioral Commitment:** The optimal value under behavioral commitment is **always no less** than that under pure commitment.



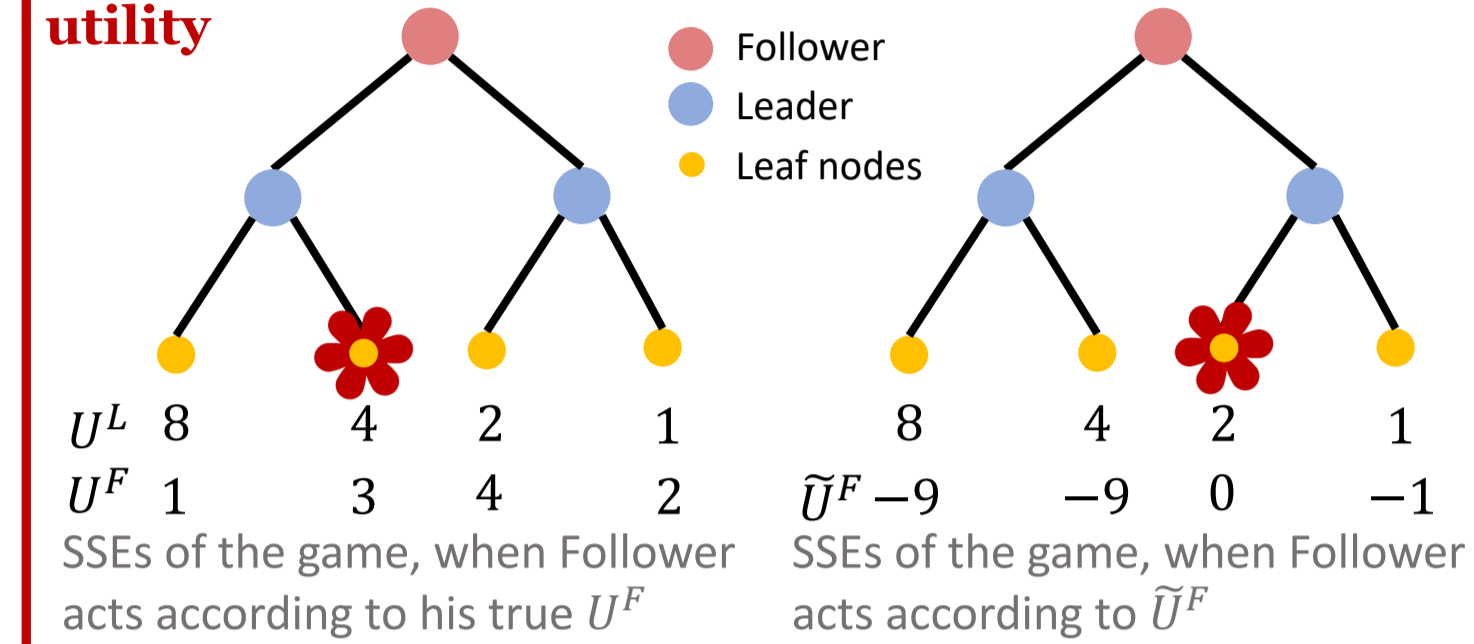
4. Pure & Inducibility

Characterization. Leaf node z is inducible iff

$$U^L(z) \geq M^L(\text{root})$$

Where $M^L(\text{root})$ is the leader's maximin value at root node.

Example: Follower deceives to gain better actual utility



Tips for deception

- show stronger conflicts of interests where the leader can gain more utilities;
- use constant-sum subgames and games with constant worst utilities to restrict the feasible strategy profiles the leader can consider to induce via commitment;

5. For More General Settings

For Pure & Strong Inducibility, Behavioral & Inducibility and Behavioral & Inducibility settings, maximin value is **NOT ENOUGH** for characterization!

- Besides the conditions related to maximin values, the same properties need to be held on subtrees.
- And other conditions are needed.

This attributes to the Tree structure, specially owned by extensive-form games.

6. A Key Technique – “Y”-shape distribution

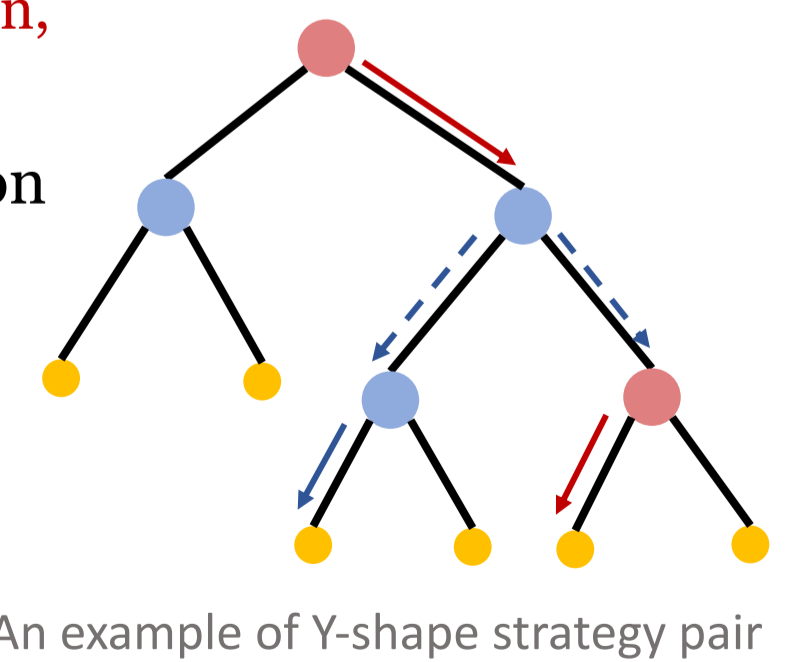
Definition. Distribution p is “Y”-shape if $\#\{\text{leaf node } z: p(z) \neq 0\} = 2$.

The corresponding strategy profiles of p yields distribution, s.t. edges with non-zero probabilities form a “Y”.

Good Property. For any (strongly-)inducible distribution p , there exists a “Y”-shape distribution p' , such that

- p' is (strongly-)inducible;
- $U^F(p') \geq U^F(p)$.

“Y”-shape distributions enable us to design algorithms for general settings and find the characterization for property USE



References:

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- Birmipas, Georgios, et al. "Optimally Deceiving a Learning Leader in Stackelberg Games." *Journal of Artificial Intelligence Research* 72 (2021): 507-531.
- Nguyen, Thanh, and Haifeng Xu. "Imitative Attacker Deception in Stackelberg Security Games." *IJCAI*. 2019.